

$$x = \{ [0, 0], [0, 1], [0, 1] \} \rightarrow \text{no remainder}$$

$$x' = [1, 0] \rightarrow \text{produce remainder}$$

$$t = [0, 0], [0, 1], [1, 1] \rightarrow \text{propagate remainder}$$

$$t' = [0, 1] \rightarrow \text{consume remainder}$$

$B \rightarrow$ some X 's

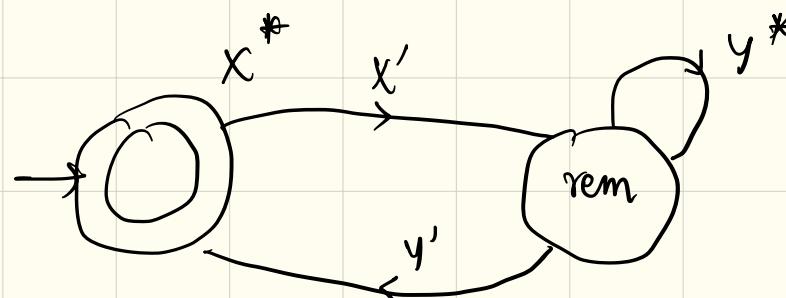
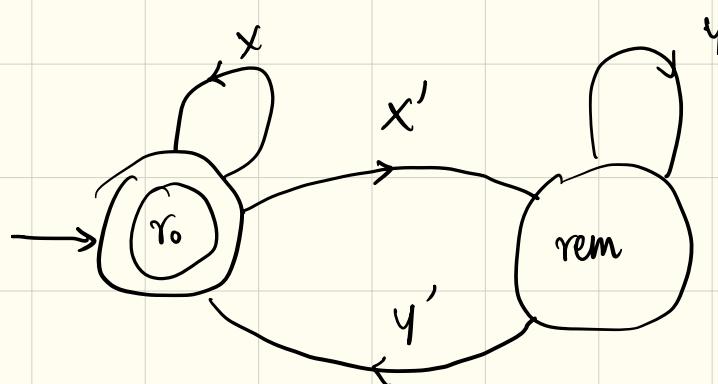
then

exactly one from X'

then some Y 's

then exactly one Y'

1 1
1 0 0
0 1 0 0
1 - 0 1
- - - -



$$\begin{bmatrix} 0 \\ 0 \end{bmatrix} \quad \begin{bmatrix} 0 \\ 1 \end{bmatrix} \quad \begin{bmatrix} 1 \\ 0 \end{bmatrix} \quad \begin{bmatrix} 1 \\ 1 \end{bmatrix} \quad 1$$

1

$$0 \rightarrow \text{remainder } 0$$

0

$$\begin{bmatrix} 0 \\ 1 \end{bmatrix} \rightarrow \begin{array}{c} 0 * 3 \\ 1 \\ 1 \\ 0 \end{array}$$

remainder 1

$$\begin{array}{ccccc} & & 3 & 0 & 1 \\ & & | & | & | \\ 1 & 1 & 3 & 0 & 1 \\ 0 & 1 & | & 1 & | \end{array}$$

1 0 0 1

1 0 1 0

0

1

1

0

$\begin{bmatrix} 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \end{bmatrix}$

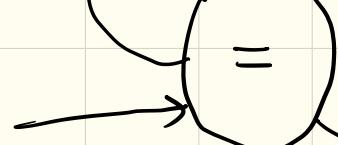
$\begin{bmatrix} 1 \\ 0 \end{bmatrix}$

1
1

0
0

0 1 0
1 0 0

$\begin{bmatrix} 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \end{bmatrix}$



Assume E is regular.

Take $p \in \mathbb{N}$ to be its pumping length

let

$$w = \begin{bmatrix} 0 \\ 1 \end{bmatrix}^p \in E$$

\downarrow
 xy^2z

s.t. $|xy| \leq p$

$\therefore y$ can only contain $\begin{bmatrix} 0 \\ 1 \end{bmatrix}$

Hence, xy^2z clearly does not belong to E .

This is a contradiction the language being regular.

$$B_n = \{ a^k \mid k \text{ is a multiple of } n \}$$

$$\begin{aligned} B_1 &= \{ a^k \mid k \text{ is a multiple of } 1 \} \\ &= a a^* \end{aligned}$$

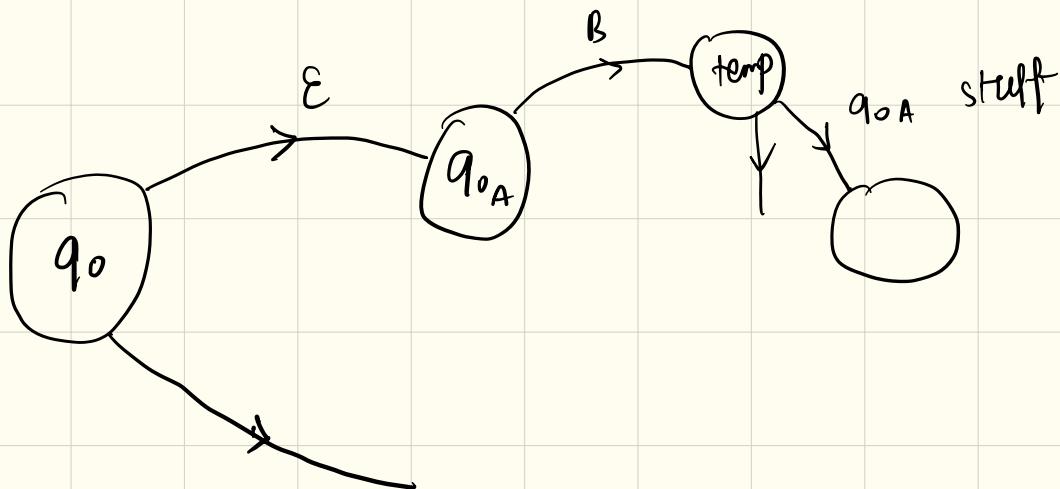
$$\begin{aligned} B_2 &= \{ a^k \mid k \text{ is even} \} \\ &= (aa)^* \end{aligned}$$

$$\begin{aligned} B_3 &= \{ a^k \mid k \text{ is multiple of } 3 \} \\ &= (aaa)^* \end{aligned}$$

1.37.

Cn

$a_1 \ b_1 \ a_2 \ b_2 \ a_3 \ b_3 \ \dots \ a_k \ b_k$



let $M_A = (Q_A, \overset{k}{\sum}, \delta_A, q_A, F_A)$

$M_B = (Q_B, \sum, \delta_B, q_B, F_B)$

$M_s = (Q_A \cup \{t_1, \dots, t_k\}, \delta', q_0, F'_A)$

$$\delta'(q, \alpha) = \begin{cases} t_i & \text{if } q = q_i \\ \delta(q_i, \alpha) & \text{if } q = t_i \end{cases}$$

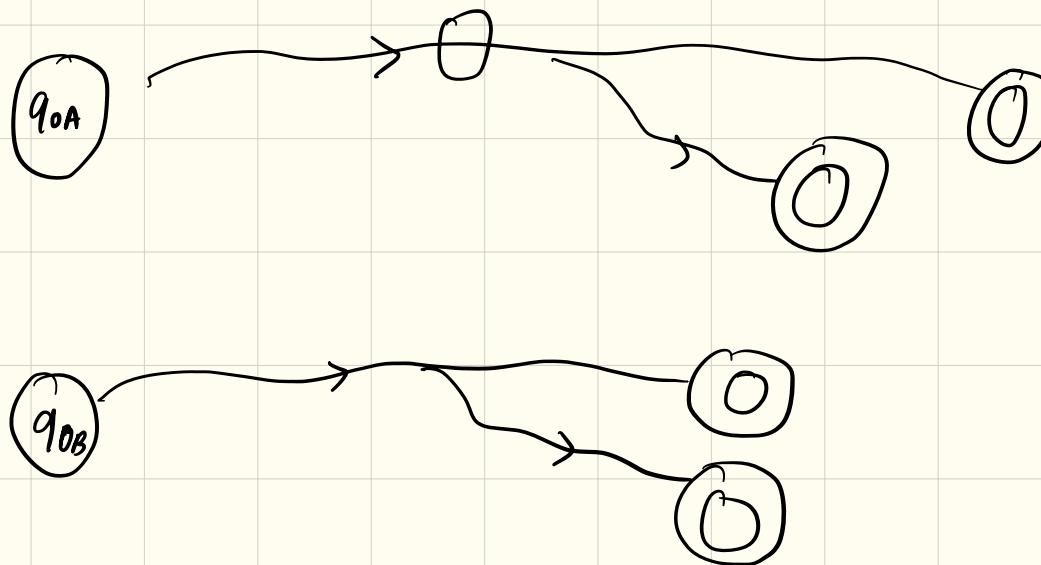
Let $a = a_1 \dots a_k$ be any string from A

Then by def^r of acceptance

$$\hat{\delta}_{M_s}(q_0, a) \in F$$

$$a' = a_1 c_1 \dots a_k c_k , \quad c_i \in \Sigma$$

$$\hat{\delta}_{M'}(q'_0, a') \in F$$



jump between machines

M_1

$$M = (Q_1 \times Q_2, \Sigma, \delta, (q_1, q_2), F_1 \times F_2)$$

NFA

$$\delta((q_1, q_2), \alpha) = \{(\delta_1(q_1, \alpha), q_2), (q_1, \delta_2(q_2, \alpha))\}$$

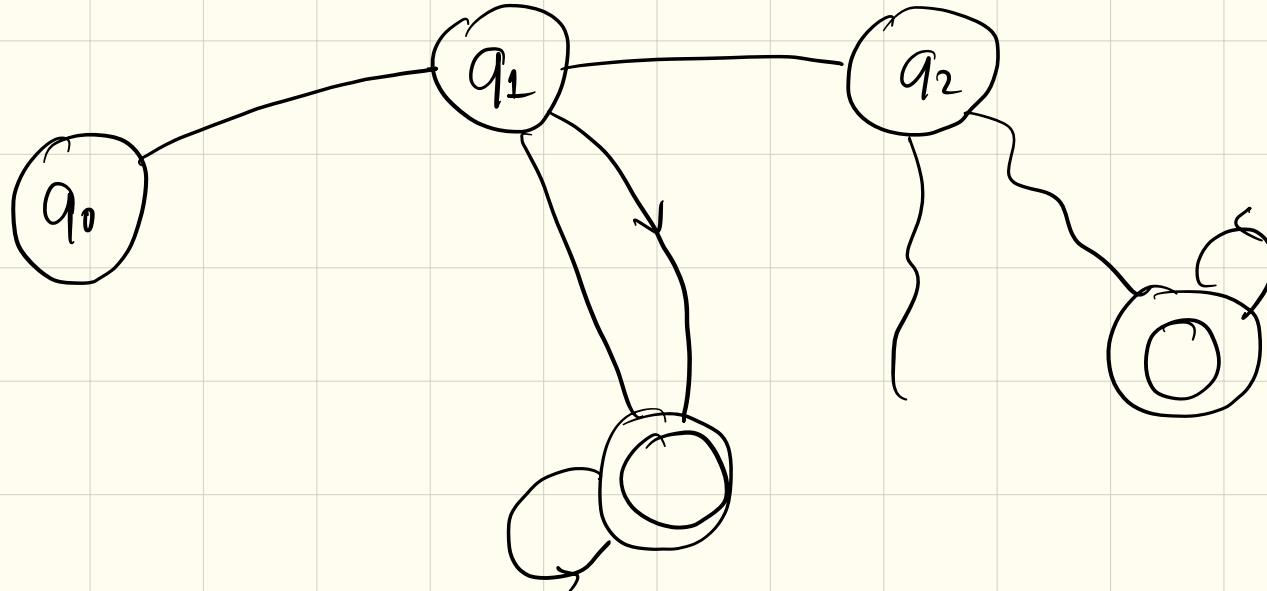
Proof: Let $a_1 a_2 \dots a_k$ be any string from A ,
then by defⁿ of M_1 $a_i \in \Sigma^*$

exists sequence of states

$$r_0, r_1, \dots, r_k \in Q_2$$

such that $r_0 = q_1, r_k \in F_1$

and for each $i = 0 \dots k-1 : \delta_1(r_i, a$



Copy the whole
 machine which accepts ,
 A
 add arrows
 from old machine
 to copied one.

$$M' = (Q \cup Q', \Sigma, \delta', q_0, F')$$

$$\delta'(t, \alpha) = \begin{cases} \left\{ \bigcup_{\beta \in \Sigma} (\delta(\delta(q, \beta), \alpha)) \right\} \cup \{\delta(q, \alpha)\} & t = q \\ \{(\delta(q, \alpha))'\}, \text{ if } t = q' \end{cases}$$

$w \in B$ s.t. for some $y \in C$

strings w and y contain equal no. of 1's

B and C parallel

usual
↓

NFA → on input 1 → same as C

on input 0 → go to all states
that can be
reached with
consecutive 0s.

$$M_1 = (Q_1, \Sigma, \delta_1, q_1, F_1) \rightarrow B$$

$$M_2 = (Q_2, \Sigma, \delta_2, q_2, F_2) \rightarrow C$$

$$M = (Q_1 \times Q_2, \Sigma, \delta', (q_1, q_2), F_1 \times F_2)$$

$$\delta((r, s), \alpha) = \begin{cases} U & \alpha = 0 \\ (\delta_1(r, \alpha), \delta_2(s, \alpha)) & \alpha = 1 \end{cases}$$

$$w \mid wx \in A \quad \text{and} \quad x \in B$$

0^P 1 0^P

0^m 1ⁿ suppose regular
→ 0^m 1ⁿ 1 0

w | w ∈ {0, 1}^{*} is a palindrome

0^P 1^P 1^P 0^P

Y = {w | w = x₁ # x₂ # ... # x_k}

$$1^p \# 1^{p-1} \# 1^{p-2} \# \dots \# 1^2 \# 1$$

y must be 1^k

1 0 1 0 0 1

1 0 1 1

0 1

$1^k y \mid y \in \{0, 1\}^*$ y contains at least k 1s

1 0 1 1 1

0^P

1 Σ^* 1 Σ^*

$$x \equiv_L x$$

$$x \equiv_L y \Rightarrow y \equiv_L x$$

$$\Rightarrow A z \in \Sigma^*$$

$$xz \in L \Leftrightarrow yz \in L$$

$$x \equiv_L y \quad y \equiv_L z$$

$$a b^n c^n$$

$$a^i$$

0 0 1

0^* 1⁺ 0⁺ 1* \cup 10 * 1

1 1 cannot be pumped

1 0 1

0 1 0

01

1 0 0

1 0 1 0 0

$$1^n$$


$$\begin{matrix} 1, & 2, & 4, & 8, & 16, & 32, \\ \downarrow & \downarrow & & \downarrow & \downarrow & \\ 1, & 2, & 11, & 21, & 121, & 1012 \end{matrix}$$

$$1 \times 3^2 + 3 \times 2 + 1$$

$$1^n \quad 2^n$$

$$1 \times 3^3 + 3 \times 1 + 2$$

$$2^n \quad 1^n$$

$$1 \ 0 \ 12$$

2 1

7

2 2 1 1

$$1 + 3 + 18 + 54$$

22

$$\begin{array}{r} 94 \\ \hline 76 \end{array}$$

1ⁿ 2ⁿ

1 2

1 1 2 2

5

$$a = 2^n - 1$$

$$b = 2^m - 1$$

$$\begin{aligned} b-a &= 2^m - 2^n \\ \overbrace{T} &= 2^n (2^{m-n} - 1) \end{aligned}$$

divisible by 2^n

$$\begin{aligned} axyz - xyz &= 3^{|y|+|z|} \cdot x + 3^{|z|} (3^{|y|} \cdot + 1) y \\ &\quad + \cancel{x} \\ &\quad - 3^{|y|+|z|} x - \cancel{3^{|z|} y} - \cancel{x} \end{aligned}$$

$$= \underbrace{3^{|y|+|z|}}_{\text{not dive}} \left(\underbrace{(3^{|y|} - 1)n}_{\text{should be divisible}} + y \right)$$

by $2^n = a+1$

But $a+1 = xyz + 1 > xyz$

$$> xyz - n$$

A regular

$$B = \{ x \mid \exists y, |y| = |x|^2 \text{ and } xy \in A \}$$

$$N = (Q, \Sigma, \delta, q, F)$$

$$N_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$$

$$Q_1 = Q_1 \times \mathcal{P}(Q)$$

$$= (q, S \mid q \in Q, S \in \mathcal{P}(Q))$$

$$q_2 = (q, F)$$

$$\delta_1((q, A), \alpha)$$

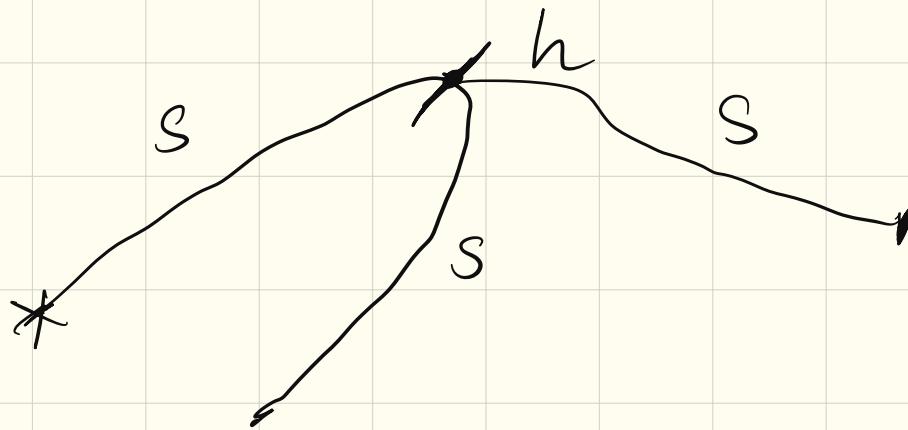
$$= (\delta(q, \alpha), \{ p \in Q \mid \exists a \in \Sigma, b \in S, \\ \delta(p, a) = b \})$$

$$F_2 = (q, S), \quad q \in S$$

$$\{ 1^n \chi 1^n \mid n \in \mathbb{N} \} \rightarrow \text{not regular}$$

\hookdownarrow contains 0

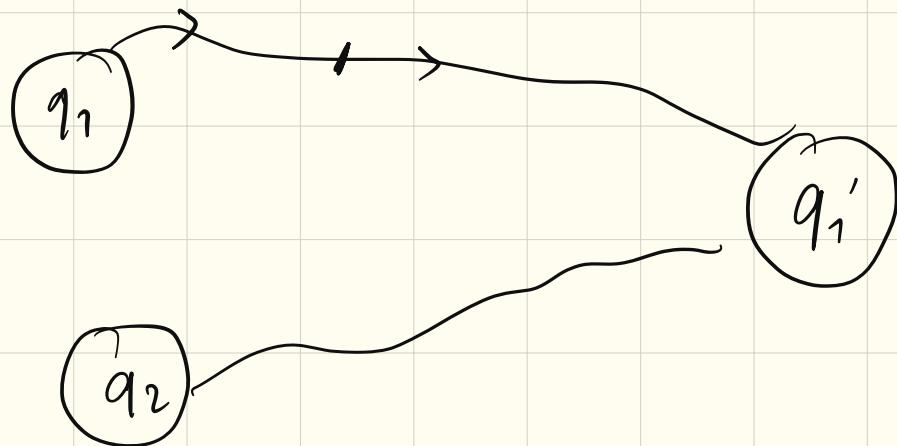
$$(1^* 0)^* 1^*$$

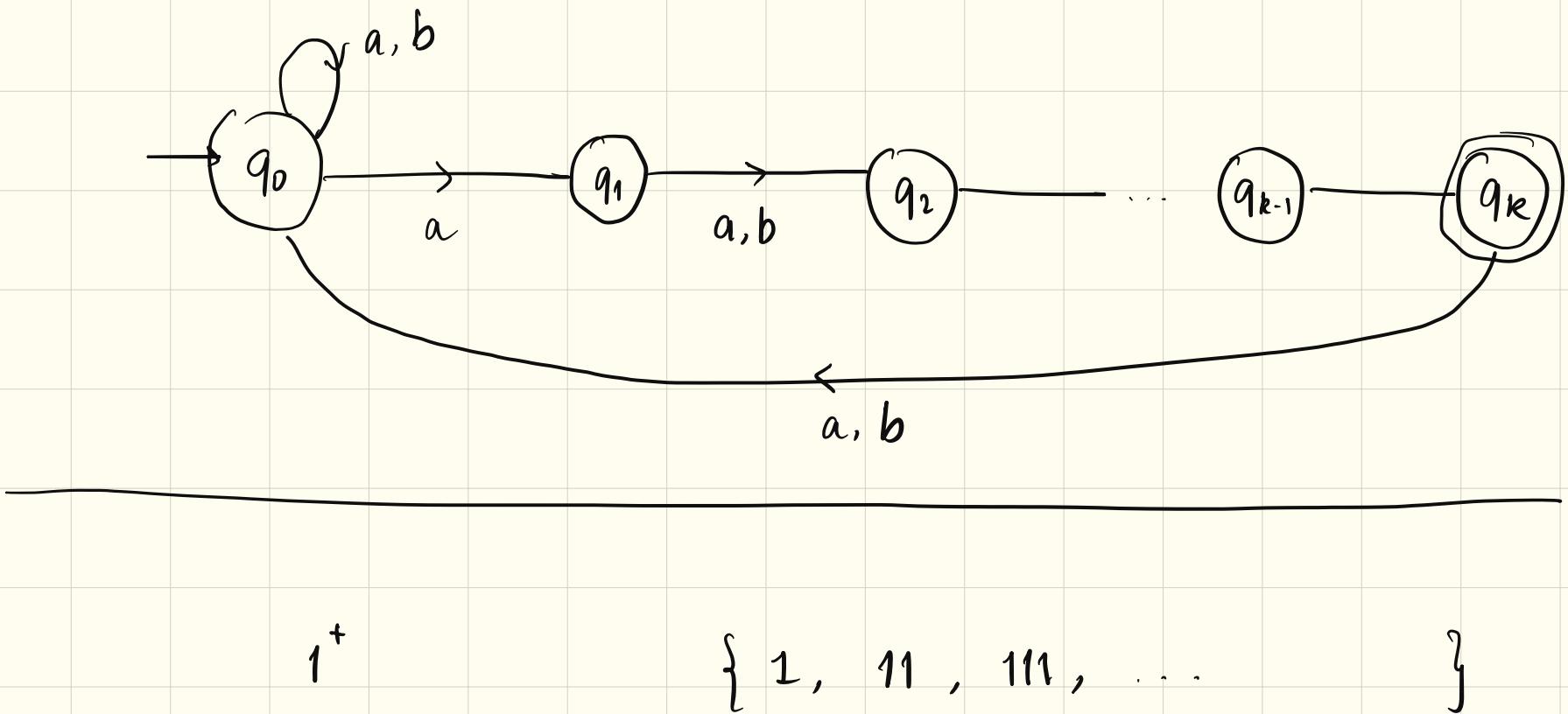


Arbitrary q_1, q_2 in a k -state synchronizable

\exists a string s s.t.

$$\hat{\delta}(q_1, s) = \hat{\delta}(q_2, s).$$



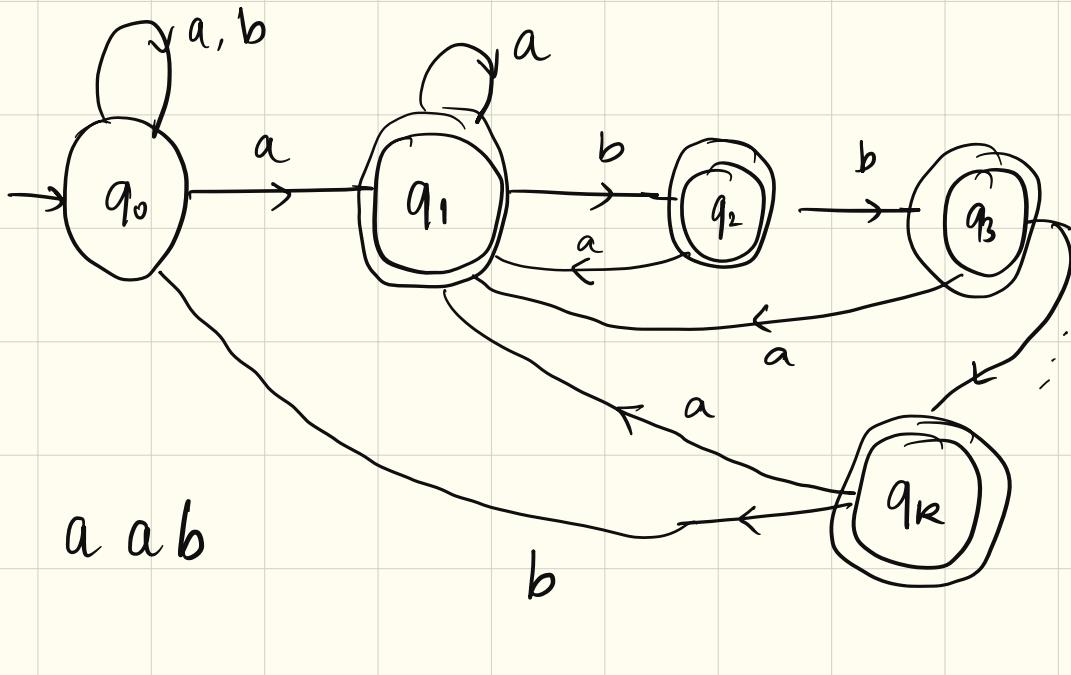


Prove by Myhill - Nerode Theorem that index of C_k is $2^k \Rightarrow$ Minimal no. of states required to recognize $C_k = 2^k$

Consider $X = \sum^k$, $\sum = \{a, b\}$ $|X| = 2^k$

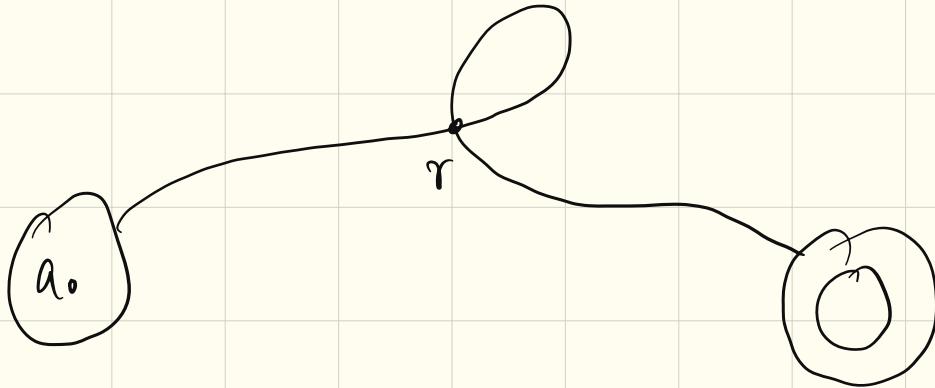
Claim: X is pairwise distinguishable by C_k (\Rightarrow index $\geq 2^k$)

Take any two distinct $x, y \in X$, then they will differ at some position say at m . One will have



$$\delta(r, \alpha) = \left\{ \begin{array}{ll} q_0 & r = q_0, \alpha = a \text{ or } b \\ q_1 & r = q_0, \alpha = a \\ q_{n+1} & r = q_n \quad \alpha = b \\ & n = 1 \dots k-1 \end{array} \right.$$

Disc → Blue
 Reddit → Orange
 YouTube → Red
 Google
 ChatGPT → cyan
 Perplexity → green
 Wikipedia
 GitHub → purple
 Gmail
 Classroom



$A_e = \{ s \mid s \in A, \text{ passes through } r \text{ even no. of times} \}$

$A_o = \{ s \mid s \in A, \text{ passes through } r \text{ odd no. of times} \}$

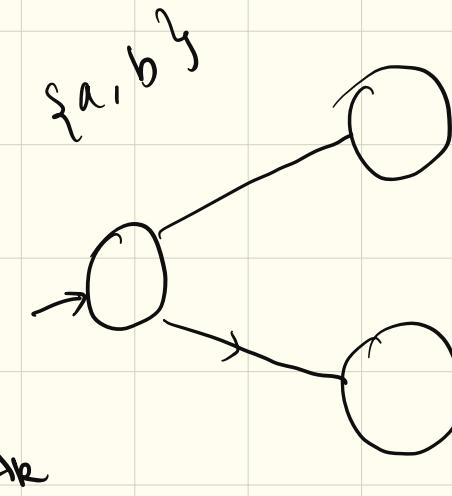
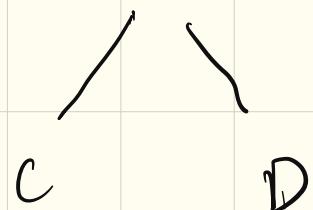
$D_{\text{even}} = (Q_e, \Sigma, \delta_e, q, f_e)$

$Q_e = \{ p, q_{oe}, q_{oo}, q_{ie}, q_{io}, \dots, q_{ke}, q_{ko} \}$

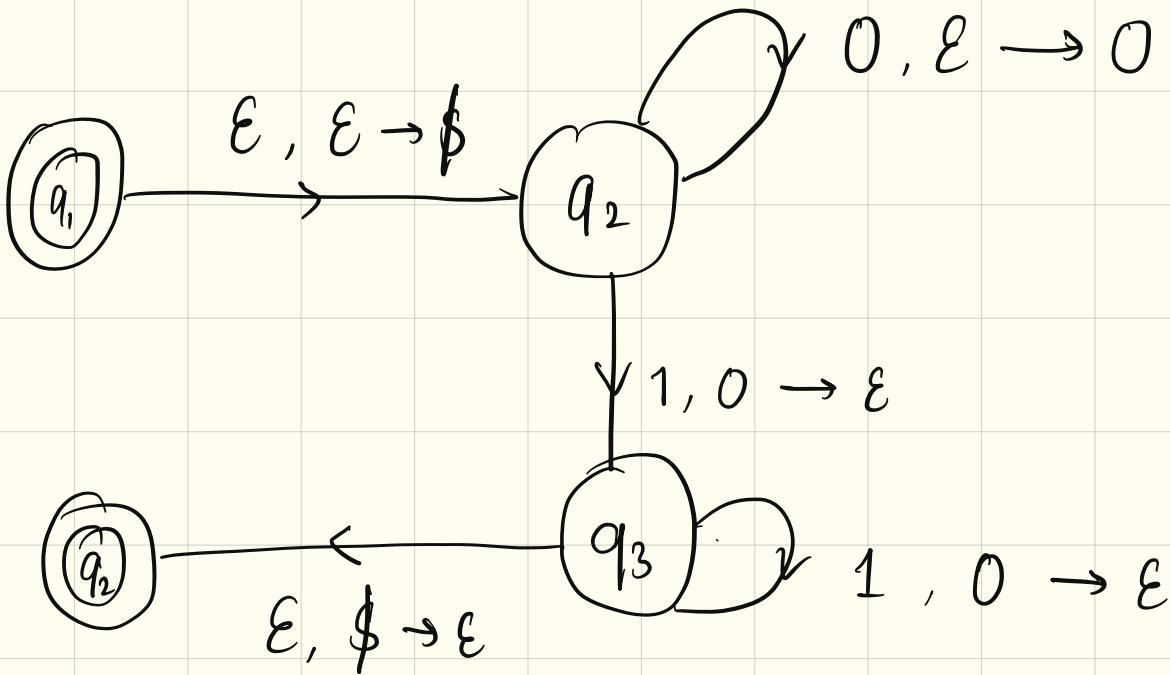
A has two disjoint regular subsets

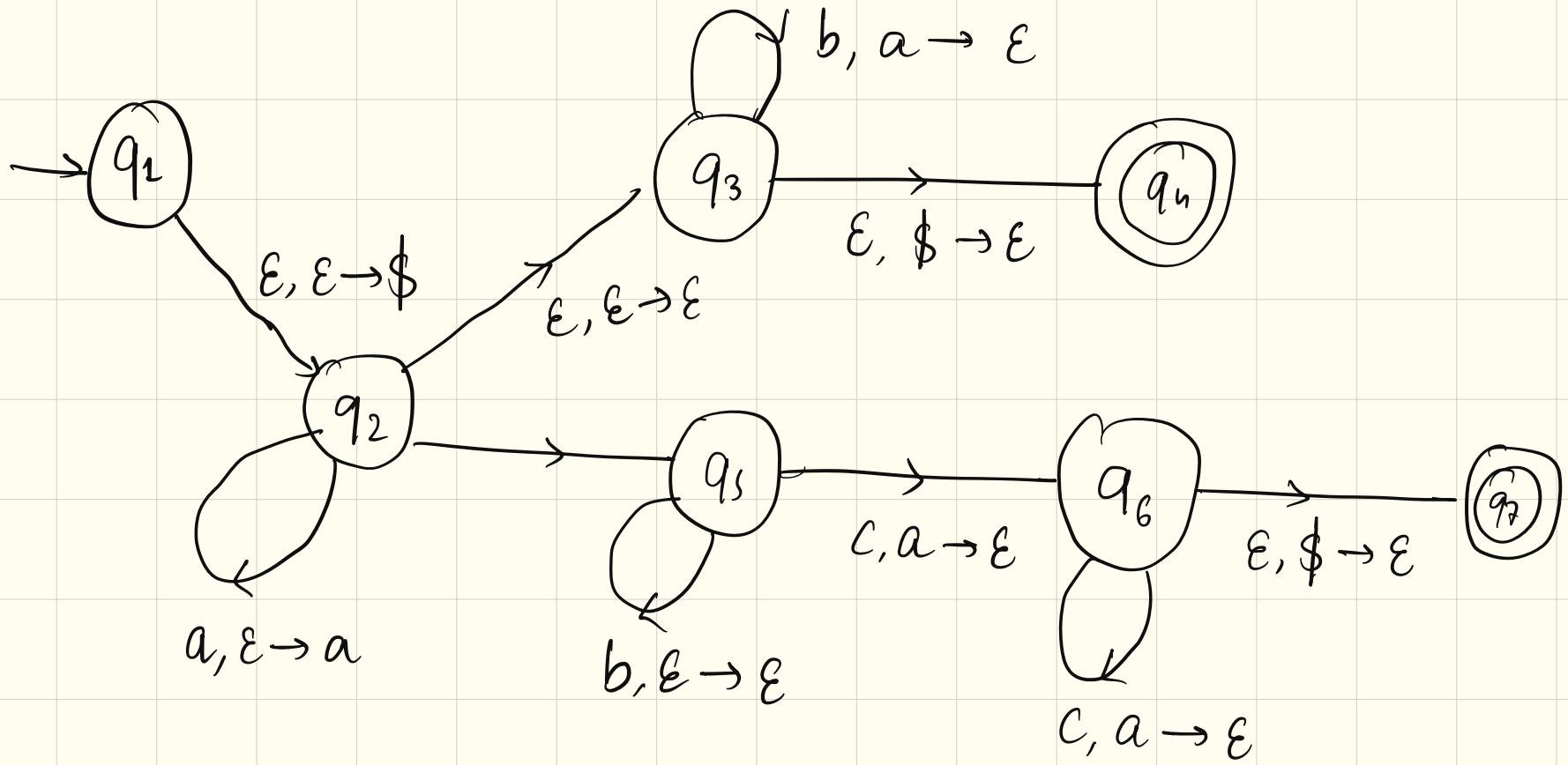
$B \subseteq D$

$D \setminus B$ is ∞



A_k

0^n 1^n 



$$S = u \vee w \vee x \vee y$$

$S \rightarrow TS \mid 0 \mid 1$

$T \rightarrow 00 \mid 01 \mid 10 \mid 11$

$S \rightarrow x0y \mid 0$

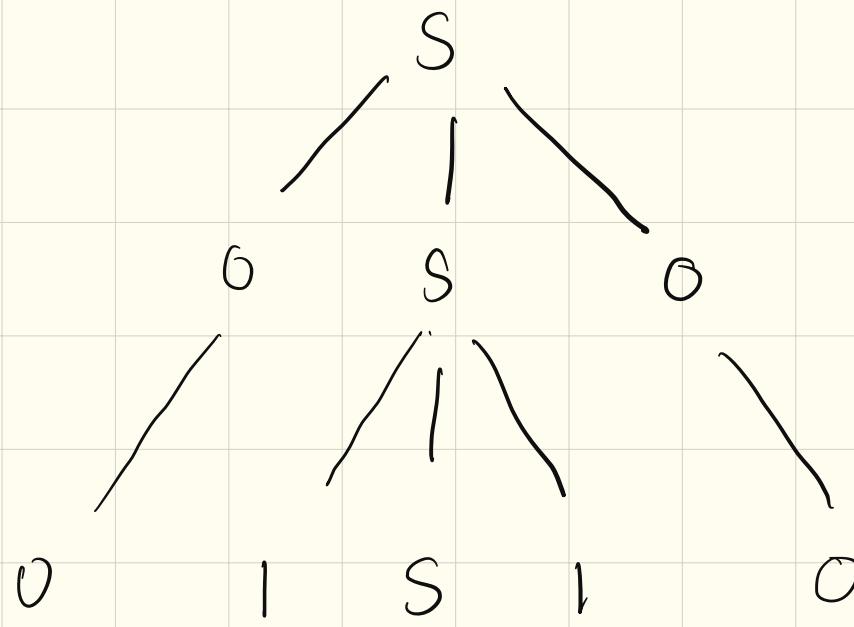
$x \rightarrow$

$x+y = \text{even}$

$S \rightarrow 0S0 | 0S1 | 1S0 | 1S1 | 0$

$\sim 1 \sim$

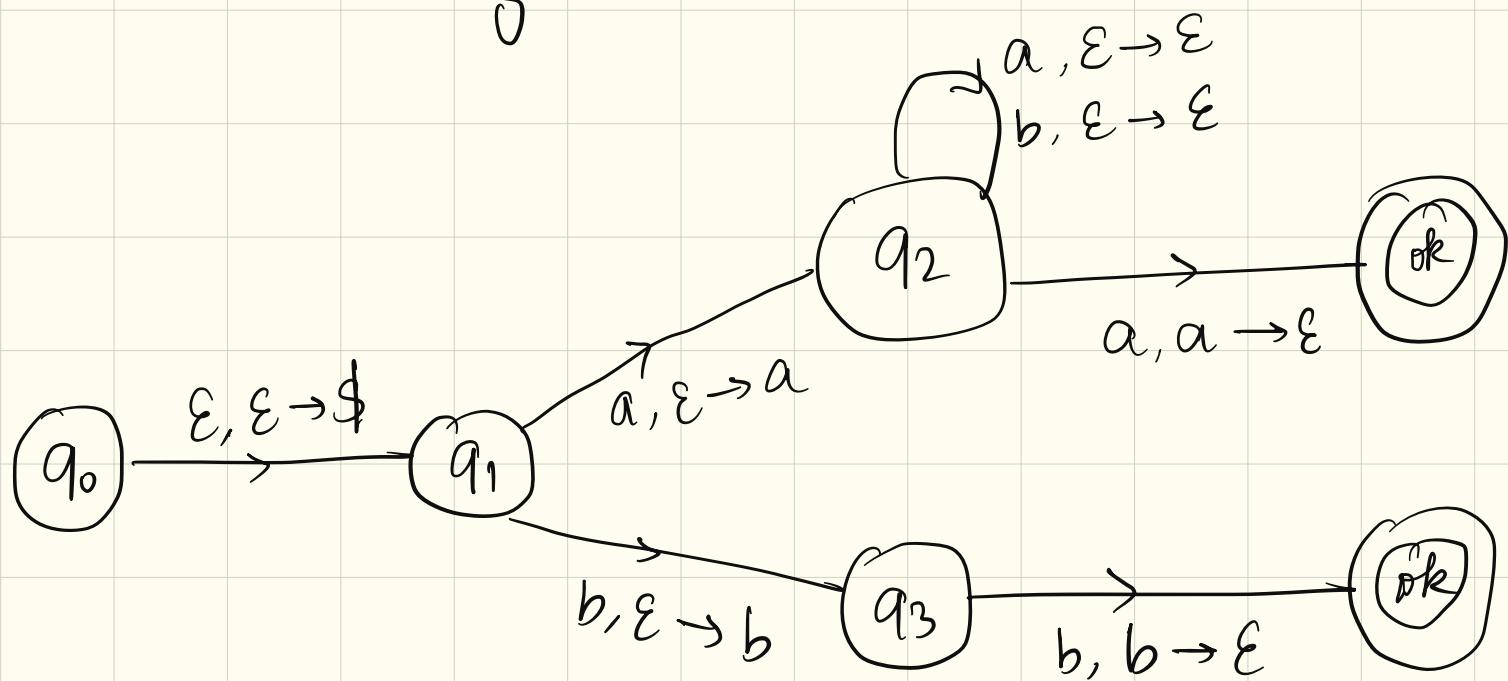
$S \rightarrow 0S0 | 1S1 | 0 | 1$

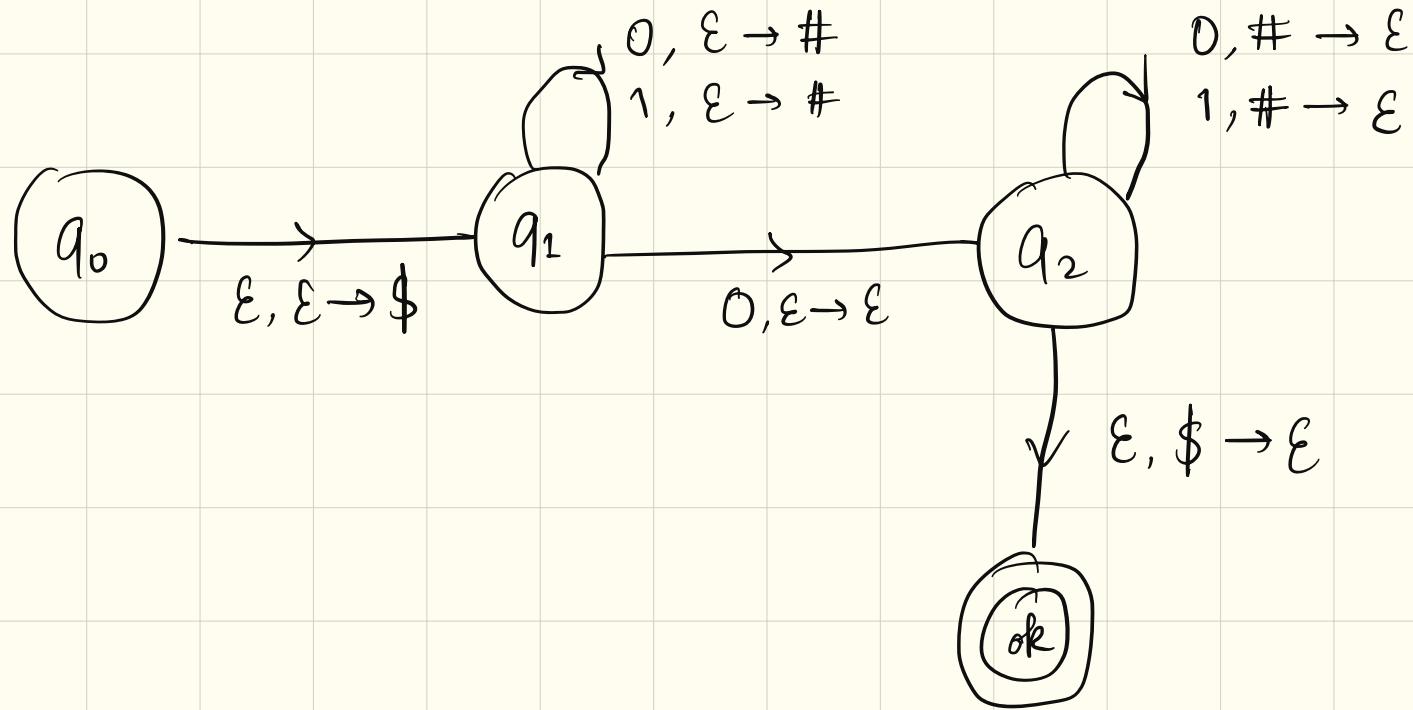


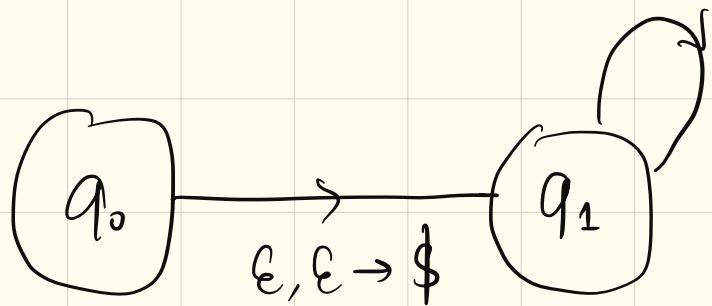
0 | | S | | 0



0



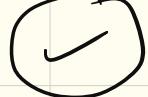




$a^n \quad b^n$

$S \rightarrow S a b \mid S b a \mid b S a \mid a S b \mid$
 $\quad \quad \quad b a S \mid a b S \mid a S \mid$
 $\quad \quad \quad S a \mid a$

$$S \rightarrow S_1 \mid S_2 \mid S_3$$
$$S_1 \rightarrow S_1 ab \mid S_1 ba \mid a S_1 b \mid b S_1 a$$
$$ab S_1 \mid ba S_1 \mid b S_1 \\ Sb \mid b$$

S_2 

$$S_3 \rightarrow Tab \mid bTa \mid baT$$
$$T \rightarrow Tab \mid Tba \mid aTb \mid bTa \\ abT \mid baT \mid \epsilon$$

$R \rightarrow Sx$

$w \# (0 \cup 1)^* w^R$

$S \rightarrow Tx$

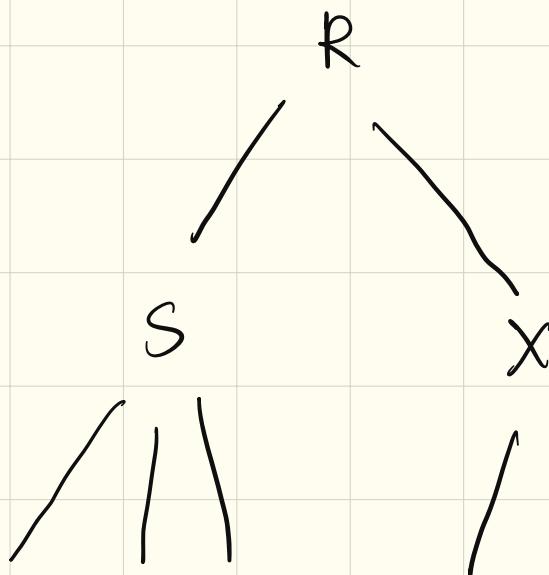
$T \rightarrow 0T0 \mid 1T1 \mid \#x$

$x \rightarrow$

0 0 #1 0 0

$S \rightarrow TX$

$T \rightarrow 0 S D 1$



0 S 0



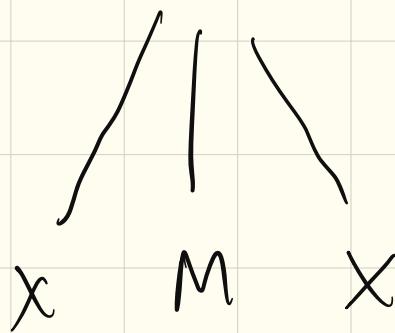
X

$$S \rightarrow LTR$$
$$L \rightarrow \epsilon \mid XM$$
$$R \rightarrow \epsilon \mid MX$$
$$M \rightarrow \# \mid \# XMX \#$$
$$T \rightarrow 0T0 \mid 1T1 \mid OM0 \mid IMI \mid \epsilon$$

0 T 0

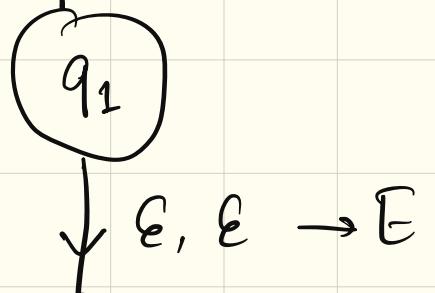
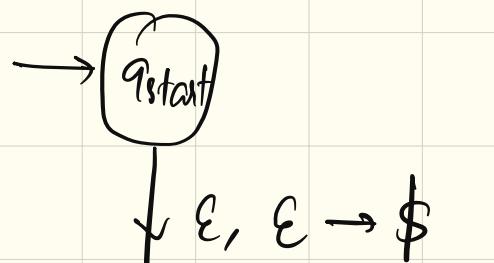


0 1 M 1 0



$$S \rightarrow S_1 | S_2$$
$$S_1 \rightarrow AS_1B | C | \epsilon$$

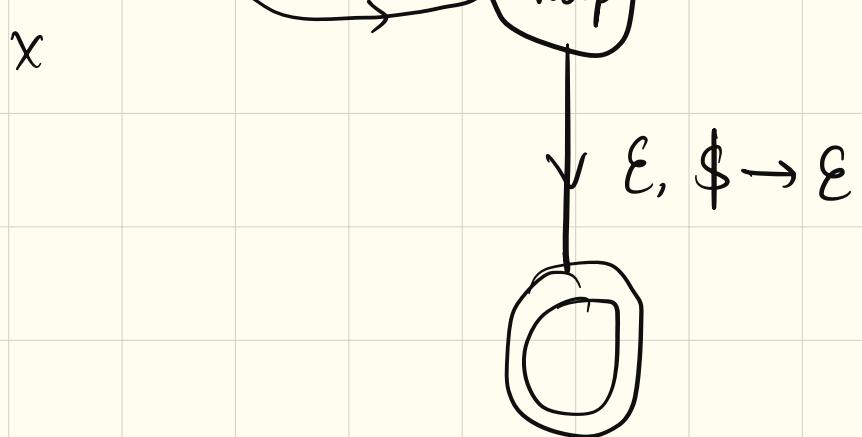
$a^n b^n$

$$C \rightarrow cC | \epsilon$$


$a, a \rightarrow \epsilon$

$t, t \rightarrow \epsilon$

q_{loop}



$$A \rightarrow BAB \mid B \mid \epsilon$$

$$B \rightarrow 00$$

$$A \rightarrow BAB \mid B \mid \epsilon$$

$$B \rightarrow CC$$

$$C \rightarrow 0$$

$$A_0 \rightarrow A \mid \epsilon$$

$$A \rightarrow BAB \mid B \mid \epsilon$$

$$B \rightarrow CC$$

$C \rightarrow O$

$A_0 \rightarrow A | \epsilon$

$A \rightarrow BAB | CC$

$B \rightarrow CC$

$C \rightarrow O$

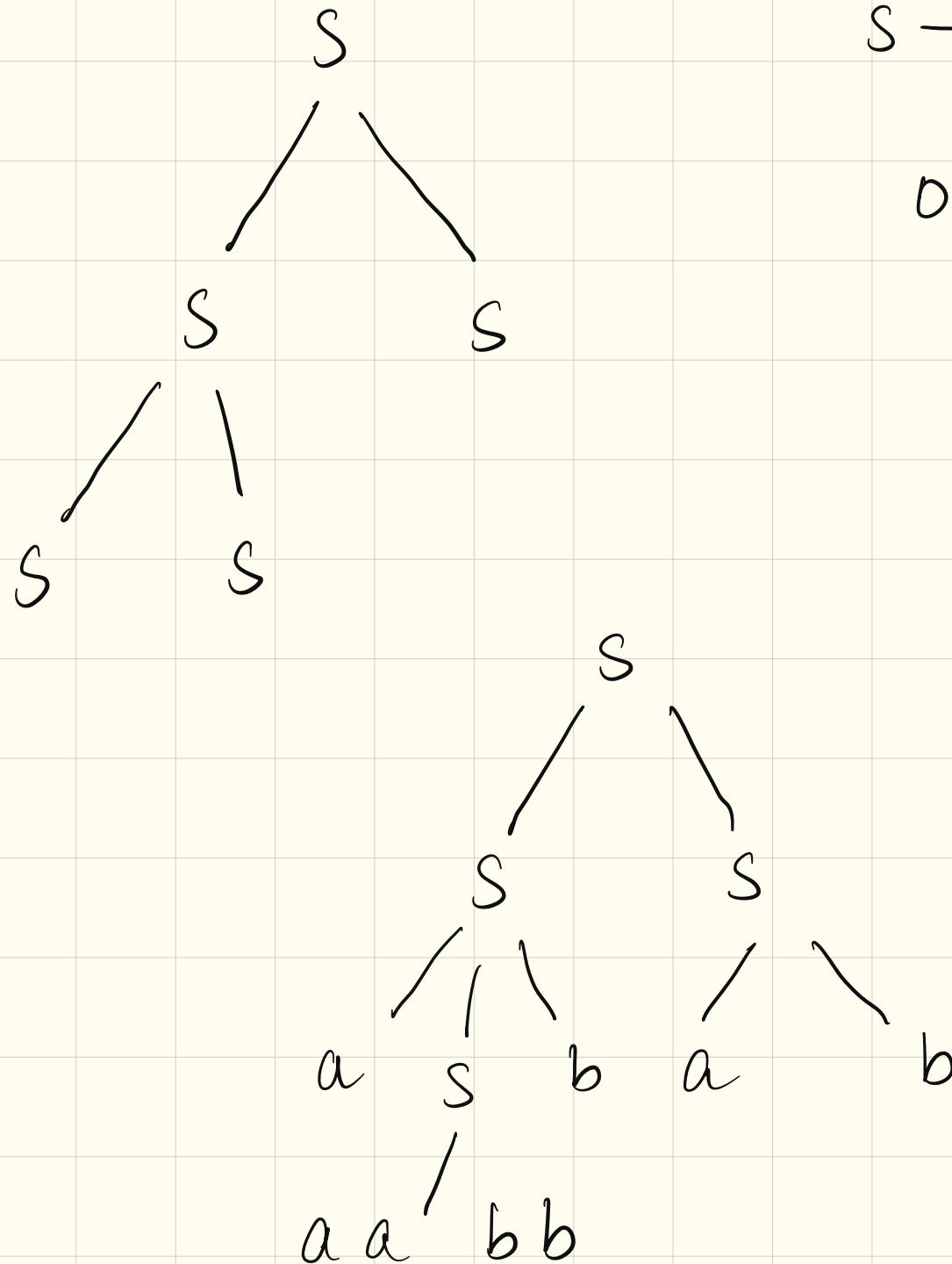
$A \rightarrow BA_1 | CC$

$A_1 \rightarrow AB$

$A_2 \rightarrow BA$

$s \rightarrow 0/1$

$0 \cup 1$



$$M = (Q_1, \Sigma, \Gamma, \delta_1, q_1, F_1) \xrightarrow{c}$$

$$N = (Q_2, \Sigma, \delta_2, q_2, F_2)$$

$$P = (Q_1 \times Q_2, \Sigma, \Gamma, \delta, (q_1, q_2), F_1 \times F_2)$$

$$S \rightarrow A \# B \mid B \# A$$

$$B \rightarrow TBT \mid 1$$

$$A \rightarrow TAT \mid 0$$

$$T \rightarrow 0 \mid 1$$

$$S \rightarrow AB \mid BC$$
$$A \rightarrow BA \mid a$$
$$B \rightarrow CC \mid b$$
$$C \rightarrow AB \mid a$$
