

25 Mar 2025 - Operating Systems - II - Week 11

Banker's Algorithm.

↳  $O(mn^2)$

BFS →  $O(n^2)$

Example of Banker's Algorithm

Total :	A	B	C
	10	5	7

Avail	A	B	C
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Allocation :

A	B	C
0	1	0
2	0	0
3	0	2
2	1	1
0	0	2

Max :

A	B	C
7	5	3
3	3	2
9	0	2
2	2	2
4	3	3

Total :

A	B	C
10	5	7

Avail

A	B	C
3	3	2

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10	5
-2	-1
-3	-1
-2	

Need:

	A	B	C
P <sub>0</sub>	7	4	3
P <sub>1</sub>	1	3	2
P <sub>2</sub>	6	0	0
P <sub>3</sub>	0	1	1
P <sub>4</sub>	4	3	1

$$\text{Need}[i] = \text{Max}[i] - \text{Allocation}[i]$$

Avail <sub>0</sub>	A	B	C	Work
	3	3	2	

$$\text{Need}_i \leq \text{avail}_i$$

P<sub>0</sub> x

P<sub>1</sub> ✓ →

P<sub>2</sub> x

P<sub>3</sub> x

P<sub>4</sub> x

Work<sub>1</sub>  
↓  
1st iteration

A	B	C
5	3	2
7	4	3

→ Allocation +

Devising this algorithm  
is complicated

↓  
applying is easy

P<sub>1</sub> ✓

Work<sub>2</sub>

P<sub>3</sub>

P<sub>4</sub>

→ now all can be satisfied  
(work ↑)  
more than req.

P<sub>s</sub> :  
↓  
single central process maintains this st.



Drawback  
SMP

every process needs to inform P<sub>s</sub> → P<sub>s</sub> becomes slow

AMP  
??  
one core

If not  $P_s$ , who?  $\rightarrow$  OS X  
 $\rightarrow$  itself a collection of processes

Obtain lock over meta lock

Since this is one single lock it will not lead to deadlock

Max  
Allocation  
Need

drawback  
lock needs to be acquired

$\rightarrow$  implement this  
compare each.

Little book on Semaphores

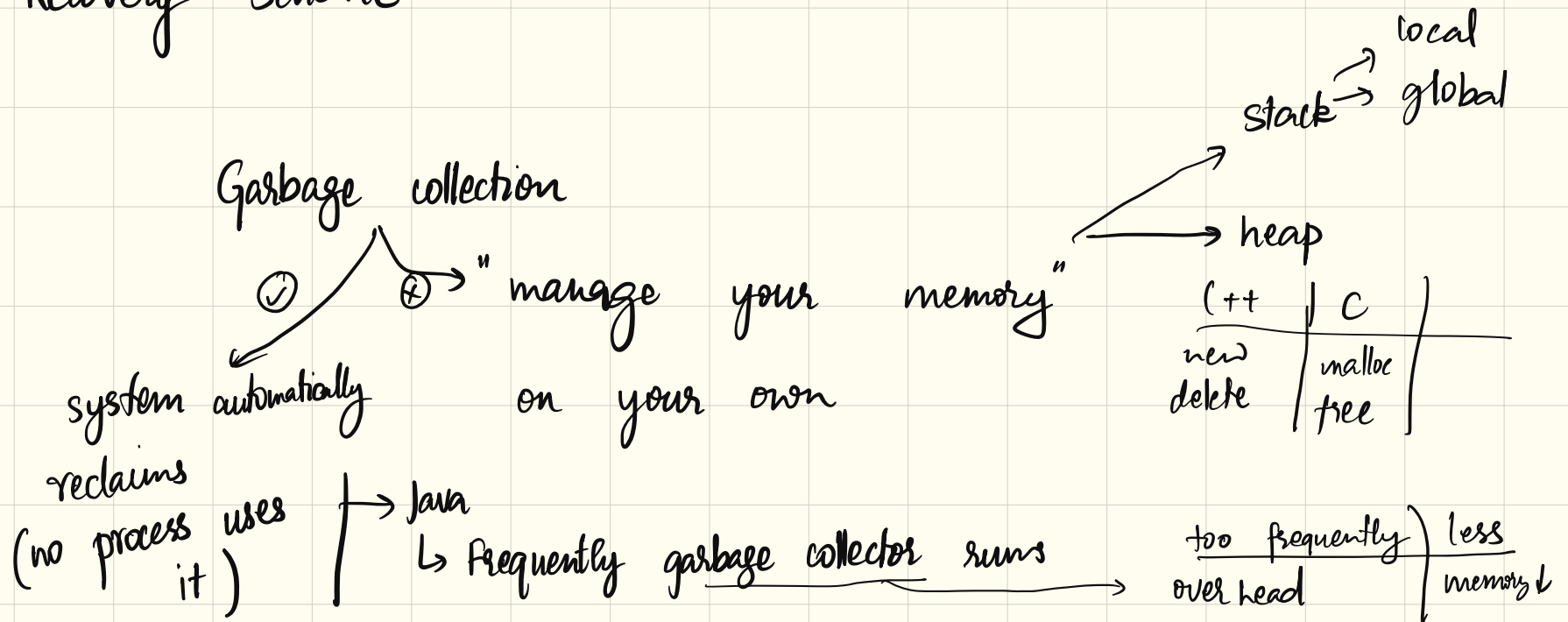
→ Not practical.

## Deadlock Detection

Allow system to enter deadlock state

Detection algorithm

Recovery scheme



Deadlock  $\rightsquigarrow$  break deadlock  
 $\searrow$  hold and wait

Resource - Allocation Graph  
Similar Algorithm  $O(mn^2)$

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→ Resource allocation graph and corresponding wait-for graph.

→ Here you don't have max no. of resources

\* Available

\* Allocation

\* S

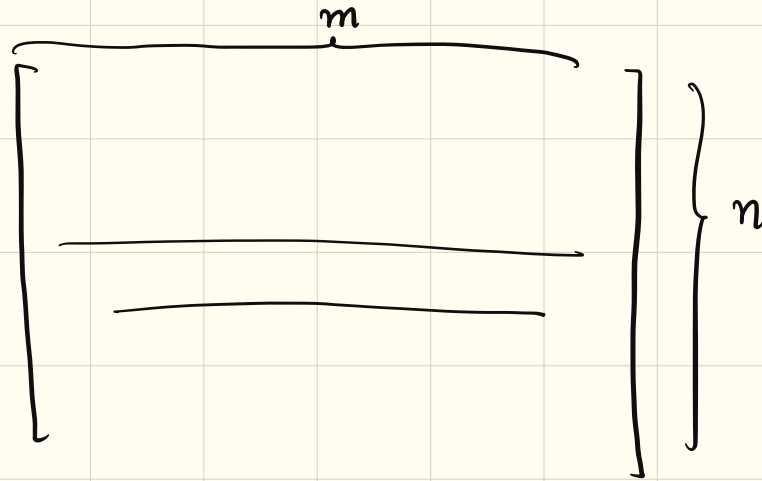
Allocation  $[i, j] = k$

$T_i$  has been allocated  $k$  instances of resource type  $R_j$

→ Detection Algorithm



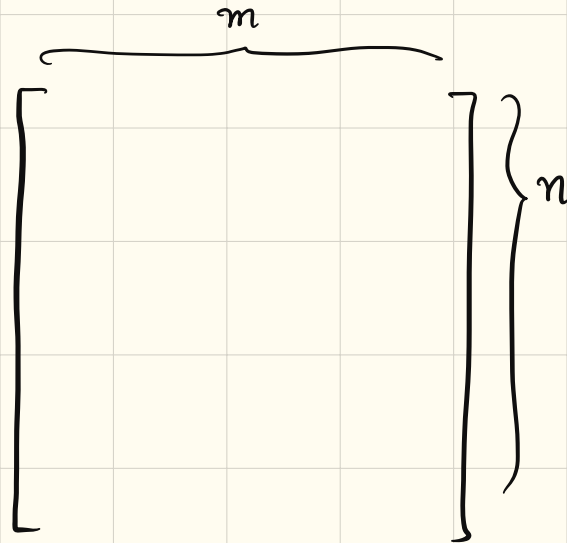
Alloc



$n \times m$

if  $Alloc_i = 0 \rightarrow$  will never be involved in deadlock  
(hold and wait)  
not satisfied

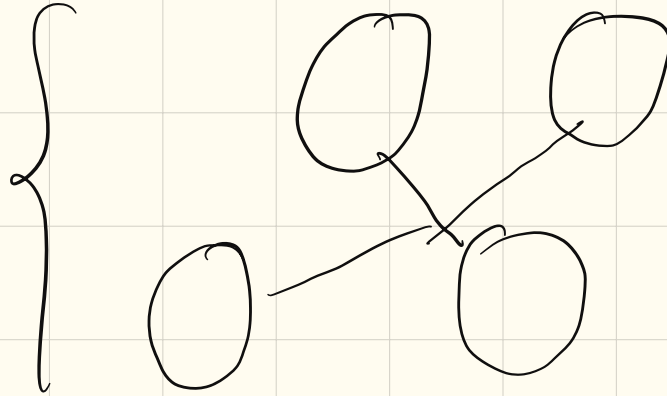
Req



$n \times m$

→ Keep checking periodically

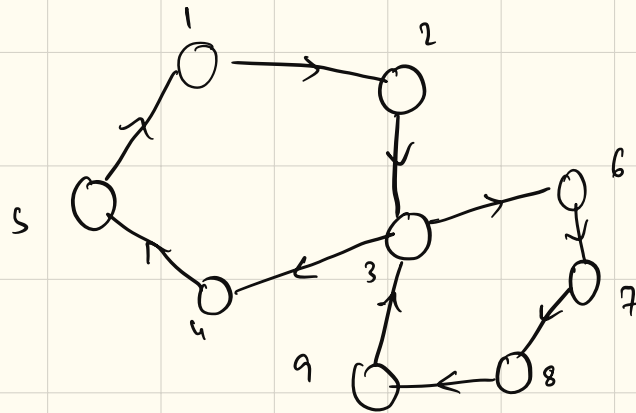
cycles in  
the graph



## Recovery from Deadlock

→ terminate a process that is involved in most no.  
of cycles

NP-complete?



- Abort and rollback
- logs for recovery
- RAID

Most popular method for deadlock prevention :

→ lock resources in order.

→ Transaction Memory

Resource Preemption

— see slides -