

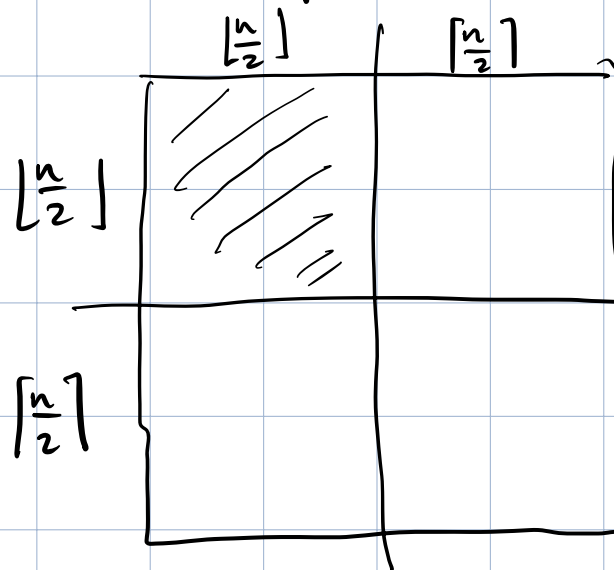
# 17 Feb 2025 - Algorithms - Week 08

Exam problems

3. Reduce to finding  $\lfloor \frac{n}{2} \rfloor^{\text{th}}$  smallest element

5. AAAH This was so easy 😞

6. Apply divide and conquer and reduce into solving in one of the 4 quads



$$T(n) = T\left(\frac{n}{2}\right) + O(n)$$

## 7. b. Radix sort

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2.  $SCS(i, j)$  = length of shortest common supersequence  
of  $A[1..i]$  &  $B[1..j]$

$$\begin{aligned} \overset{i=1}{SCS(1, j)} &= j && \text{if } A[1] \in \{B[1], \dots, B[j]\} \\ &= j+1 && \text{o/w} \end{aligned}$$

### Recurrence

$A[1] \quad \dots \quad A[i]$

$B[1] \quad \dots \quad B[j]$

If  $A[i] = B[j]$

$$SCS(i, j) = SCS(i-1, j-1) + 1$$

If  $A[i] \neq B[j]$

$$SCS(i, j) = 1 + \min(SCS(i-1, j), SCS(i, j-1))$$

4.  $P(j, i) = \Pr$  [ India wins atleast  $j$  games in first  $i$  ]  
 $0 \leq j \leq i$

$$\Pr(j, i) = p_j \cdot \Pr(j-1, i-1) + \Pr(j, i-1) \cdot (1 - p_j)$$

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$P(j, i) = \Pr$  [ India wins exactly  $j$  games in first  $i$  ]  
 $0 \leq j \leq i$

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Running time :  $O(n^2)$

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for j = 0 to n
  for i = j to n
    find P(i, j)
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Edit distance

Running time :  $O(mn)$

Space :  $O(mn)$

can improve to  $O(m+n)$

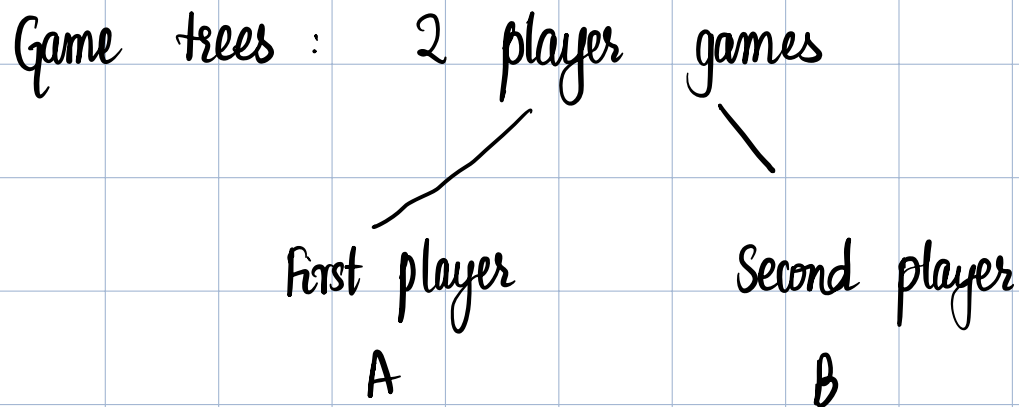
How to backtrack?

??

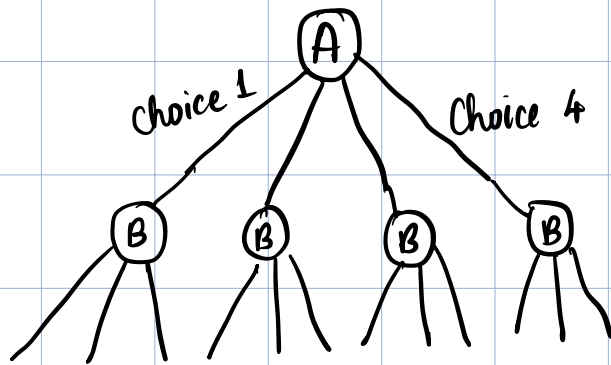
Graph  
markup  
language

# Dynamic Programming on Trees

Input = tree, instead of a sequence



→ no chance involved in the outcome of each move.



Leaf nodes represent end of game.

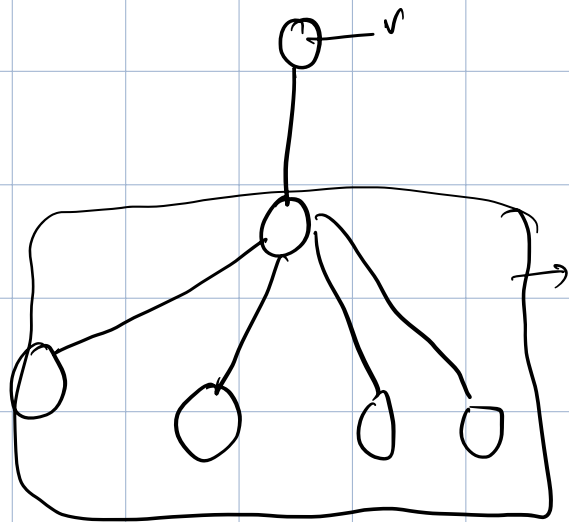
→ value represents who won the game.

Every finite 2-player game with no ties has a winning strategy for one of the players.

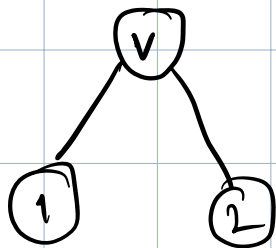
A is the first player → makes move first.

Possible Questions: Which player has a winning strategy and what is it?

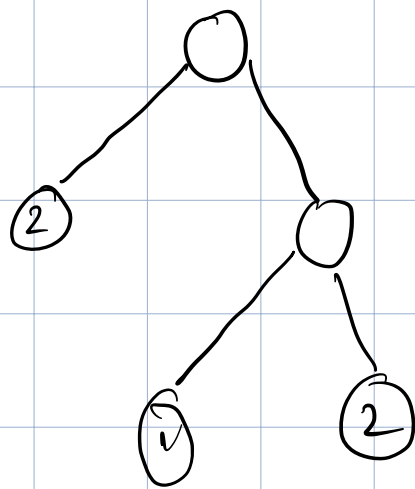
Consider an arbitrary subtree



$$f(v) = 1 \quad \text{if player at } v \text{ wins } T_v \\ = 0 \quad \text{otherwise}$$



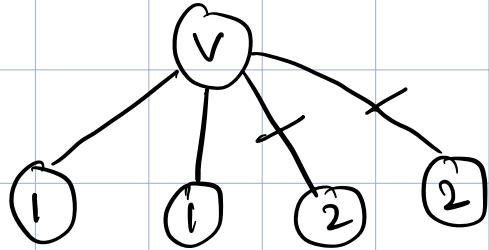
player 1 wins



} 2 always wins

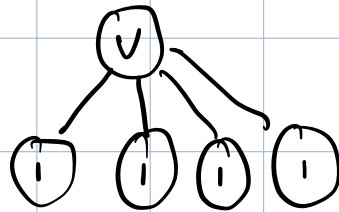
Compute  $f(v)$  for  $v =$  leaves then parents of  
leaves and so on.





v can win

$$f(v) = 1$$



v cannot win

$$f(v) = 1$$

$$f(v) = 2 \quad \text{if}$$

$$f(w) = 1 \quad \forall \text{ children } w \text{ of } v$$

A's turn at  $v$  :

$f(v) = A$  if at least one child  $w$  of  
B has  $f(w) = A$

$f(v) = B$  if at least one child has  
 $f(w) = B$

